CS 173
Tandy Warnow
Topics for today

• Basics of combinatorial counting
• Applications to running time analysis
Using combinatorial counting

- Evaluating exhaustive search strategies:
  - Finding maximum clique
  - Determining if a graph has a 3-coloring
  - Finding a maximum matching in a graph
  - Determining if a graph has a Hamiltonian cycle or an Eulerian graph
Combinatorial counting

How many ways can you
• put $n$ items in a row?
• pick $k$ items out of $n$?
• pick subsets of a set of size $n$?
• assign $k$ colors to the vertices of a graph?
• match up $n$ boys and $n$ girls?
Technique

• To count the number of objects, design an algorithm to generate the entire set of objects. Check if each object is created exactly once (if not, you will have to do a correction later).

• The algorithm’s output can be seen as the leaves of a decision tree, and you can just count the leaves.
Putting n items in a row

Algorithm for generating all the possibilities:
• For i=1 up to n, DO
  – Pick an item from S to go in position i
  – Delete that item from the set S

Analysis: each way of completing this generates a different list.

The number of ways of performing this algorithm is n!
Number of subsets of a set of size \( n \)

Algorithm to generate the subsets of a set

\[ S = \{s_1, s_2, s_3, \ldots, s_n\} \]

- For \( i=1 \) up to \( n \) DO:
  - Decide if you will include \( s_i \)

Analysis: each subset is generated exactly once, and the number of ways to apply the algorithm is \( 2 \times 2 \times \ldots \times 2 = 2^n \).
k-coloring a graph

Let $G$ have vertices $v_1, v_2, v_3, \ldots, v_n$

Algorithm to k-color the vertices:
• For $i=1$ up to $n$ DO:
  – Pick a color for vertex $v_i$

Analysis: each coloring is produced exactly once, and there are $k^n$ ways of applying the algorithm.
Matching n boys and girls

Algorithm:

• Let the boys be $B_1, B_2, \ldots B_n$ and let the girls be $G_1, G_2, \ldots G_n$.

• For i=1 up to n DO
  – Pick a girl for boy $B_i$, and remove her from the set

Analysis: there are n ways to pick the first girl, n-1 ways to pick the second girl, etc., and each way produces a different matching.

Total: $n!$
Picking k items out of n

Algorithm for generating all the possibilities:
• For i=1 up to k, DO
  – Pick an item from S to include in set A
  – Delete that item from the set S

The number of ways of performing this algorithm is \( n(n-1)(n-2)\ldots(n-k+1)=n!/(n-k)! \)

But each set A can be generated in multiple ways - and we have overcounted!
Fixing the overcounting

Each set A of k elements is obtained through k! ways of running the algorithm. As an example, we can generate \{s_1, s_5, s_3\} in 6 ways, depending upon the order in which we pick each of the three elements.

So the number of different sets is the number of ways of running the algorithm, divided by k!.

The solution is \( \frac{n!}{[k!(n-k)!]} \)
Summary (so far)

• To count the number of objects, design an algorithm to generate the entire set of objects. Check if each object is created exactly once (if not, you will have to do a correction later).

• The algorithm’s output can be seen as the leaves of a decision tree, and you can just count the leaves.
Summary

• Number of orderings of n elements is n!
• Number of subsets of n elements is $2^n$
• Number of k-subsets of n elements is $n!/\left[k!(n-k)!\right]$
• Number of k-colorings of a graph is $k^n$
More advanced counting

• What is the number of k-subsets of a set $S = \{s_1, s_2, s_3, \ldots, s_n\}$ that do not include $s_1$?

• What is the number of k-subsets of a set $S = \{s_1, s_2, s_3, \ldots, s_n\}$ that do include $s_1$?

• What is the number of orderings of the set $S$ in which $s_1$ and $s_2$ are not adjacent?

• What is the number of orderings of the set $S$ in which $s_1$ and $s_2$ are adjacent?
New techniques

- Count the complement
- Divide into disjoint cases, and count each case
Example

• What is the number of k-subsets of a set $S = \{s_1, s_2, s_3, \ldots, s_n\}$ that do not include $s_1$?
• Solution: same as number of k-subsets of $\{s_2, s_3, \ldots, s_n\}$. So $(n-1)!/[(n-1-k)!k!]$
Example

• What is the number of k-subsets of a set $S = \{s_1, s_2, s_3, \ldots, s_n\}$ that do include $s_1$?

• Solution: same as the number of $(k-1)$-subsets of $\{s_2, s_3, \ldots, s_n\}$, so $(n-1)!/[(n-k)!(k-1)!]$
Example

• What is the number of orderings of the set $S$ in which $s_1$ and $s_2$ are adjacent?

• Solution: two cases:
  – Case 1) $s_1$ followed by $s_2$
  – Case 2) $s_2$ followed by $s_1$

• Same number of each type. Easy to see that there are $(n-1)!$ of each type, so $2(n-1)!$ in total
Example

• Number of orderings of the set $S$ in which $s_1$ and $s_2$ are not adjacent?

• This is the same as $n!-2(n-1)!$
Using combinatorial counting

• Evaluating exhaustive search strategies:
  – Finding maximum clique
  – Determining if a graph has a 3-coloring
  – Finding a maximum matching in a graph
  – Determining if a graph has a Hamiltonian cycle or an Eulerian tour